Tracking Network Flows with P4

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What is P4?

- P4: Programming Protocol-Independent Packet Processors
- Programming language that describes the behavior of the data plane
- Allows the rapid development and deployment of protocols
- An alternative to fixed function devices (e.g. Switches, Routers, Firewalls…)
- P4\textsubscript{14} with software switch provided by the P4 Language Consortium
Motivation

Can the features provided by P4 be used to enhance the ability to identify and track malicious traffic in networks?

- **Goals**
  - Identify as much of the path of a flow as possible
  - Minimize impact on system resources

- **Assumptions**
  - Traffic may be classified as malicious at later time
  - NetFlow is unable to capture all flows at every node

Sampled NetFlow may result in incomplete knowledge of the path.
Path Tracking Methods: Hop Recording

- Each node adds its node ID to a header stack
- Flow and Path saved at egress node
- Not a novel approach
  - RFC 791 - Internet Protocol, 1981
  - In-band Network Telemetry (INT), 2015
- Variable amount of data added to packet
- Useful when load balancing over several short paths
Path Tracking Methods: Forwarding State Logging

- Network forwarding state identifier added to packet
- Use a record of forwarding state to recreate path
- Requires global view of how routing will be performed
- Requires archives of each version of routing database
Path Tracking Methods: Dynamic Path Labeling

- MPLS style labels are generated for each path taken by packets
- Labels identify the history of a packet rather than future
- Final label stored with flow at egress node
Bloom Filters

- Probabilistic Set
- Possible false positives, No false negatives
- Configurable Accuracy
- Fixed size data structure
- Independent of size of item being added
How Bloom Filters Work

- Bloom Filter is made up of fixed number of bits ($m$)
- Has predetermined number of hash functions ($k$)
- The hash functions determine which bits would be set if an item was in the set
- Items can only be added, not removed
- As items are added, chance for a false positive increases
Bloom Filters in P4

- Can be implemented completely in data plane using P4
- Used to determine when a new flow has been seen
- Control Plane will only receive packets that are from a new flow
- Data Plane can send to centralized point without involving the Control Plane

```p4
bit bloom_filter[BF_WIDTH]

action bf_check_bit(salt) {
    index = hash(hash_fields, salt) % BF_WIDTH
    bit = bloom_filter[index]
    hit = hit & bit
}

action bf_set_bit(salt) {
    index = hash(hash_fields, salt) % BF_WIDTH
    bloom_filter[index] = 1
}
```
 Practical Example

- Set a goal of 64000 flows \( (n) \) with a false positive rate of 5\% \( (p) \)
- 64001st flow has a 5\% chance of falsely being detected as previously seen
- Actual false positive rate for the first 64000 flows was 1.16\%
Counting Bloom Filters

- Can now tell how many times a flow has been seen
- Uses substantially more memory
- Possible concurrency issues
- Inaccuracies now result in inflated counts
Decrementing Bloom Filters

- Counteract a Bloom Filter ‘filling up’
- Can decrement based on time or packets
Application of Counting Bloom Filters

- Detecting SYN attacks using ratio of SYN to non-SYN packets
- Relatively low-bandwidth attack
- Assuming single or aggregate source
SYN Attack Detection Performance

- SYN attack is using 20 packets per second
- Lower decrement rate increases accuracy
- Lower performance the more distributed the attack
Conclusions

- Need to look at the additional constraints that hardware will impose
- The effectiveness of the flow tracking methods depends on the topology of the network and resource constraints
- Bloom Filters can provide high accuracy with minimal resource utilization
- What effect will moving to P4\textsubscript{16} have on implementation?
Questions?